

## Empirical Investigation of Cloud, Grid and Virtualization Using Compiler Optimization Level for Threads Processes

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### ABSTRACT

This research focused on implementation of OpenMP. It considers the parallelization of an application code which simulates the thermal gradient of a material in two dimensions. A C language program code called **jacobi2d.c** that solves a rectangular 2 dimensional heat conductivity problem using Jacobi iterative method was used. The boundary conditions required to compute a temperature distribution for a rectangular 2D problem are: Top 300C, Bottom 500C, Left 400C and Right 900C with a range of problem sizes enter as a run-time parameters to alter the problem sizes and convergence criteria. Also, there were computations and readings for iterations and runtime for four values of M and N which was selected for 01, 02 and 03 optimizations. In Table 1.1 Readings, four values were selected for each of the iterations. The results shows the performance of the runtime as the processor increases from 01 -optimization, to 02 -optimization and finally to 03-optimization. It can be deduce from the representation that the run time of the values reduces as more resources are allocated to execution through the increase in optimization level. Also, in Table 2.1 Readings, the runtime decreases as it moves from thread1, thread2, thread3 and thread4, comparing the last values for thread1 which are M is 180,N is 200, and their runtime which is 42.797187001 . Also the last values for thread2 which are M is 180, N is 200, their runtime which is 21.772106003. When the two runtimes were compared, it was discovered that there was a decrease in the runtime because the more the thread increases, the more system resources they share such as a processor which may affect their runtime by increasing it.

**Keywords:** Cloud, Grid, Virtualization, Threads, Processors, Compiler, Optimization.

### 1. BACKGROUND TO THE STUDY

Cloud computing emerges as a new computing paradigm which aims to provide reliable, customized and QoS guaranteed computing dynamic environments for end-users (Wang et al, 2008). A cloud computing is an approach where large scale related capability computer resources and infrastructures are provided in form of services across the Internet to numerous customers. A grid machine can be described as an infrastructure that can be used for solving dynamic problems, such as multi-processes, resources allocation and a centralized control mechanism, using a standard set of protocols and interfaces to deliver significant quality services (Cafaro, M., & Aloisio, G. 2011).

The enormous computing resources demand of a process can be solved by a parallel computing implementation specifically developed to work in Grid environments of multiprocessor computing resources. The different parallel computing approaches (intra-node, inter-node and inter-organisations) are not sufficient to address the computing resources demand of such a big problem (Aparício et al, 2006).

## 1.2. Concept Oo Cloud, Grid and Virtualization

This research enhances the context of Grids, Clouds, and Virtualization. Grids computing ensure the delivery of computing power and resources on demand. However, despite the various contributions of active research in the area of grid computing, no viable commercial grid computing provider has emerged. In daily activities, some users or consumers of computational resources will always need to make provision for their own supercomputers that can guarantee speed, timely delivery and multi-processing.

The concept of Cloud Computing is not a completely new approach, research has established that there is a relative correlation between Grid Computing, cloud computing and other relevant technologies such as utility computing, cluster computing, and distributed systems in general (Foster et al, 2008). The concept of Virtualization is a technology that enables many different

Clouds to be integrated. Some researchers focused the definition of grid, cloud and virtualization around on-demand access to computing, data, and services. A grid system comprises of hardware and software infrastructure that provides dependable, consistent, pervasive, and inexpensive access to high-end computational capabilities (Foster, I., & Kesselman, C. 2003) . The concept of Cloud computing is an information technology approach that is characterized by ubiquitous access to shared resources and services that can be provided rapidly with minimal interference, via the Internet facilities.

In cloud computing, there is sharing of resources such as software and hardware in order to achieve a coherence services and cost saving approach. In the field of computer science, a thread of execution comprises of the smallest sequence of executable instructions that can be processed independently by a scheduler, which forms part of the operating system. Different operating systems processes threads and processes differently from on another, but mostly a thread is made up of the component of a process. Many threads can executed within a single process, executing concurrently and sharing the same computer resources such as memory, while some processes do not share memory resources or computer resources. Most times, the threads of a process can share its executable code and the values of its variables at any given time (Lamport, 1979).

The issue of Virtual Machine (VM) concept dates back to the early 60s, this approach was introduced by IBM as a mean to provide concurrent, interactive access to their mainframe computers. A VM was an example of the physical machine and gave users the illusion of accessing the physical machine directly. Virtual Machine was developed and used to enable time-sharing and resource-sharing at the same time on the expensive hardware resources. Virtualization has been helpful in reducing the cost of hardware resources and to improve the overall productivity by accommodating as many users as possible to work on it simultaneously (Cafaro, M., & Aloisio, G. 2011). However, as technology advances and become more integrated, the hardware has become cheaper and affordable and at the same time multiprocessing operating systems has been invented and integrated (Chiueh, S. N. T. C., & Brook, S. 2005).

### 1.3 Statement of Problem

The development of mobile and portable computing devices has created an enabling environment for minimizing the power consumed by a computer program. In the early development, the limitations of the computer memory have been a limiting factor which hampers the performance of optimizations. Because of all these influence, optimizations do not often produce optimal result, and in fact an optimization sometimes may become an impediment to performance of resources. As a result of these problems, there is a need for compiler optimization and compares of runtime in threads processes in order to determine the processor performance.

### 1.4 Objective

The objective of this research is the Implementation of an OpenMP which considers the parallelization of an application code that simulates the thermal gradient of a material in two dimensions using C language program code called `jacobi2d.c` and compares the results of the optimization and threads processes in order to determine the performance of the runtime as the processor increases.

## 2. METHODOLOGY

The study adopted a combination of qualitative and quantitative research methods. It explores the concept of cloud, grid and virtualization. Also utilizes parallelization of an application code which simulates the thermal gradient of a material in two dimensions using a C language program code called `jacobi2d.c` that solves a rectangular 2 dimensional heat conductivity problem using Jacobi iterative method to test the runtime and determine the performance of the processor. This research analysis was conducted using University of Greenwich UK CMS grid machine resources.

### 3. DATA PRESENTATION AND ANALYSIS

#### 3.1 Compiler Optimization

##### A) Step 1

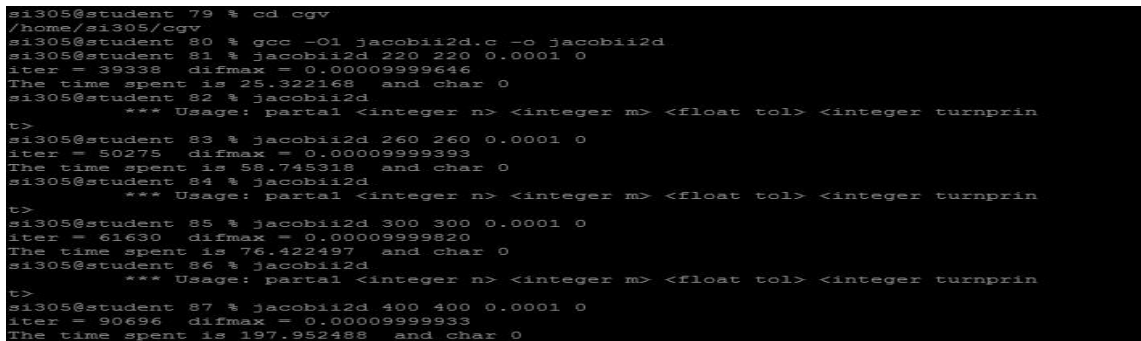
In the step one of this research, there is a modification to the jacobi2d.c code to reflect the following boundary conditions, at top 30C, bottom 50C, left 40C, and at the right 90C. The tolerance was set to 0.0001, the result was set to not printing using 0, and 1 for printing of the boundary sizes. Four different values were selected for M and N for different optimization, which comprises of 01-optimization, 02-optimization and 03-optimization.

##### Reflection of boundary sizes.

Compiler optimization is a process to minimize the time taken to execute a program. In this step 1, the jacobi2d.c code was modified to:

```
// fix boundary conditions
for (i=1; i <= m; i++) {
    t[i][0] = 40.0;
    t[i][n+1] = 90.0;
}
for (j=1; j <= n; j++) {
    t[0][j] = 30.0;
    t[m+1][j] = 50.0;
```

Fig 1: Reflection of boundary sizes.



```
si305@student 79 % cd cgv
/home/si305/cgv
si305@student 80 % gcc -O1 jacobi2d.c -o jacobi2d
si305@student 81 % jacobi2d 220 220 0.0001 0
iter = 39338 difmax = 0.00009999646
The time spent is 25.322168 and char 0
si305@student 82 % jacobi2d
*** Usage: partial <integer n> <integer m> <float tol> <integer turnprin
t>
si305@student 83 % jacobi2d 260 260 0.0001 0
iter = 50275 difmax = 0.00009999393
The time spent is 58.745318 and char 0
si305@student 84 % jacobi2d
*** Usage: partial <integer n> <integer m> <float tol> <integer turnprin
t>
si305@student 85 % jacobi2d 300 300 0.0001 0
iter = 61630 difmax = 0.00009999829
The time spent is 76.422497 and char 0
si305@student 86 % jacobi2d
*** Usage: partial <integer n> <integer m> <float tol> <integer turnprin
t>
si305@student 87 % jacobi2d 400 400 0.0001 0
iter = 90696 difmax = 0.00009999933
The time spent is 197.952488 and char 0
```

Figure 2: Screen shot for Compiler, program name and runtime parameters.

### 3.2 Optimization Level

The following computations are the readings of iterations and runtime for the four values of M and N which was selected for 01, 02 and 03 optimizations. Four values were selected for each of the iterations, and the table below shows the four values and the result of the iteration and the runtime for each of the iterations performed.

**Table 1: Readings for 01,02,03 optimization Runtime for values M and N**

**Compiler : gcc jacobii2d.c -o jacobii2d , Executable name: jacobii2d**

READINGS FOR	Value M	Value N	Runtime	Iterations	Difmax
<b>01- OPTIMIZATION</b>	220	220	25.322168	39338	0.00009999646
	260	260	58.745318	50275	0.00009999393
	300	300	76.422497	61630	0.00009999820
	400	400	197.952488	90696	0.00009999933
READINGS FOR	Value M	Value N	Runtime	Iterations	Difmax
<b>02- OPTIMIZATION</b>	220	220	16.488465	39338	0.00009999646
	260	260	28.701726	50275	0.00009999393
	300	300	46.694705	61630	0.00009999820
	400	400	151.077582	90696	0.00009999933
READINGS FOR	Value M	Value N	Runtime	Iterations	Difmax
<b>03- OPTIMIZATION</b>	220	220	16.314877	39338	0.00009999646
	260	260	27.160550	50275	0.00009999393
	300	300	45.176557	61630	0.00009999820
	400	400	130.857994	90696	0.00009999933

The table above comprises of the readings for 01, 02, 03 optimizations for four selected values of M and N. It was discovered from the readings and the results of the runtime that the runtime decreases as the optimization level increases; this indicates that the processor allocates more resources and thereby increasing the rate of execution of the runtime. In order to explain this in a more graphical form, a bar chart was used to illustrate the runtime performance. This bar chart explains the performance of the runtime as the processor increases from 01-optimization, to 02-optimization and finally to 03 -optimization. We could deduce from the graphical representation that the run time of the values reduces as more resources are allocated to execution through the increase in optimization level.

From the table above, the last runtime for each of the optimization values was selected, which comprises of 01 optimization, 02-optimization and 03-optimization.

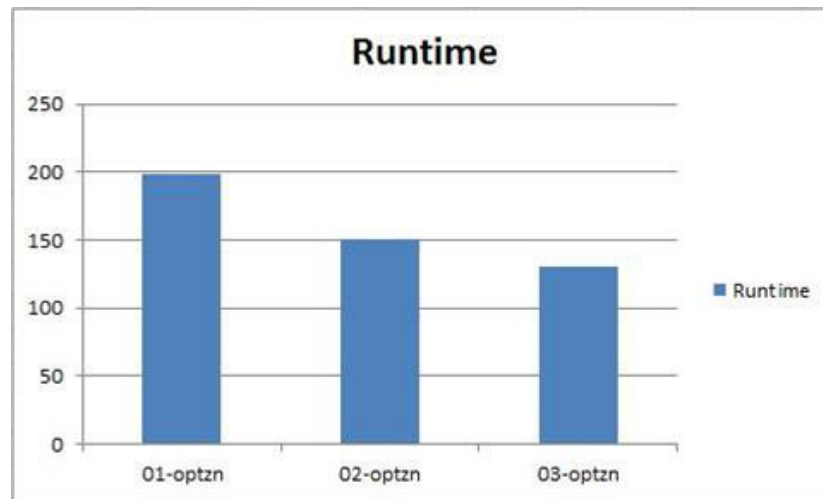
For all the last values of the each optimization:

**M is 400, N is 400**

Computing the runtime for the extracted last values in each optimization, we have:

**Table 2: Extracted last values of M, N in each optimization**

Optimizations	Runtime
01-optimization	197.952488
02-optimization	151.077582
03-optimization	130.857994



**Figure 3: Runtime Result of 01,02,03 optimizations for value M,N**

In summary of the chart analysis above, it is recorded that for values M 400,N 400 with following runtime (197.952488) for 01-optimization, runtime (151.077582) for 02-optimization and runtime(130.857994) for 03-optimization, it was discovered that when the processors are optimized, or when higher optimization are used to run a set of values, the runtime tends to decrease with respect to the increase in optimization. The result of the chart above, can therefore be concluded that the blue lines which signifies the rate of the runtime decreases with the respect to increase in optimization.

### 3.3 OpenMP PARALLEL VERSION OF Jacobi CODE

#### B) STEP 2

In step 2, there is a modification to the application created in step1 to be able to get parallel of the code with openmp using the code (`#pragma omp parallel for default (shared) private(i,j)`). Timer was also included using the timer code

(`omp_get_wtime()`;) so as to determine the parallel runtime of the code, it was also tested on four (4) threads or processors to be able to measure the performance and record the parallel runtime.

```
#pragma omp parallel for default(shared) private(i,j)

for (i=1; i <= m; i++) {

for (j=1; j <= n; j++) {

tnew[i][j] = (t[i-1][j]+t[i+1][j]+t[i][j-1]+t[i][j+1])/4.0;

#pragma omp parallel for default(shared) private(i,j,diff)

for (i=1; i <= m; i++) {

for (j=1; j <= n; j++) {

diff = fabs(tnew[i][j]-t[i][j]);

if (diff > difmax) {

difmax = diff;
```

Figure 4: Openmp Parallel cod

### 3.4 Threads Level

Table 3: Readings for 1,2,3,4 Threads Parallel Runtime for values M and N

Compiler: gcc -fopenmp jacobiopenmp.c -o jacobiopmmp, Executable name: jacobiopenmp

READINGS FOR THREAD	Value M	Value N	Runtime	Iterations	Difmax
1	70	80	1.489187002	7589	0.00009998674
	100	120	5.872882999	14161	0.00009997808
	140	160	18.256341003	23698	0.00009998894
	180	200	42.797187001	34672	0.00009999279
2	70	80	0.807676002	7589	0.00009998674
	100	120	3.106940001	14161	0.00009997808
	140	160	9.461168002	23698	0.00009998894
	180	200	21.772106003	34672	0.00009999279
3	70	80	0.742388003	7589	0.00009998674
	100	120	2.592556998	14161	0.00009997808
	140	160	7.853976000	23698	0.00009998894
	180	200	16.931387000	34672	0.00009999279
4	70	80	0.576768998	7589	0.00009998674
	100	120	1.973710999	14161	0.00009997808
	140	160	5.704780001	23698	0.00009998894
	180	200	14.126476999	34672	0.00009999279

#### 4. DISCUSSION OF FINDINGS

The table 2.1 above shows the various readings for value M and N for about four (4) threads, the readings consists of the values, runtime, iterations and Difmax for each of the value used. The same set of values was used for all the four threads in order to be able to allow comparism between the values, especially their runtime. Based on the readings taken from the four (4) threads, the following were discovered:

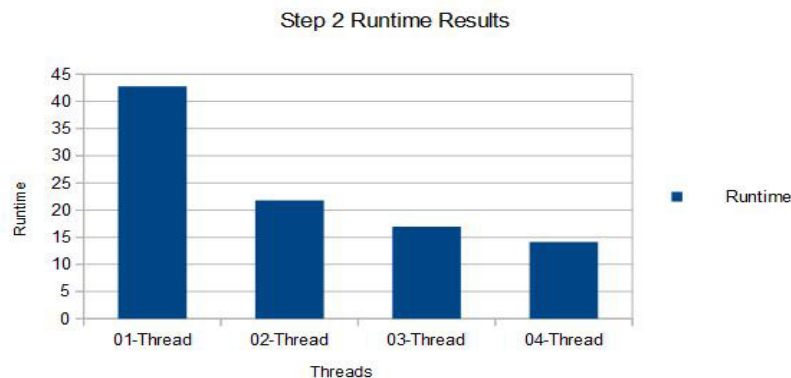
- (4) The runtime decreases as it moves from thread1, thread2, thread3 and thread4, comparing the last values for thread1 which are M is 180,N is 200, and their runtime which is 42.797187001 . Also the last values for thread2 which are M is 180, N is 200, their runtime which is 21.772106003. When the two runtimes were compared, it was discovered that there was a decrease in the runtime because the more the thread increases, the more system resources they share such as a processor which may affect their runtime by increasing it.
- (5) The iterations are the same for all the values for thread1, thread2, thread3, thread4.
- (6) The Difmax are also the same for thread1, thread2, thread3, thread4.

In addition to the results of the readings above in table 2.1, the graph below interprets the runtime for each thread. For the graph, the last values of each thread were used for thread01, thread02, thread 03 and thread04. The last values of all the threads are:

**M is 180, N is 200**

**Table 4: Extracted last values of M and N in each thread.**

01-Thread	42.797187001
02-Thread	21.772106003
03-Thread	16.931387000
04-Thread	14.126476999



**Figure 5: Step 2 Runtime results**

This graph shows additional information on the decrease in runtime as the number of threads increases.



## **5. CONCLUDING REMARKS**

This research work demonstrates using OpenMP to parallelize a practical application. It shows how parallel performance tuning using OpenMP as well as compiler optimizations can be used to achieve improved performance. The parallelization of an application code which simulates the thermal gradient of a material in two dimensions was used. The research work explains the performance of the runtime as the processor increases from 01-optimization, to 02-optimization and finally to 03-optimization. We could deduce from the result representation that the run time of the values reduces as more resources are allocated to execution through the increase in optimization level. Also, it was discovered that there was a decrease in the runtime because the more the thread increases, the more system resources they share such as a processor which may affect their runtime by increasing the run time.

## **6. CONTRIBUTIONS TO KNOWLEDGE**

Based on the study and research work conducted, and in order to demonstrate how parallel performance and compiler optimization can be used to achieve improve processor performance, the following recommendations are suggested: Better performance can be achieved when the grid machine resources is not been freeze or overload by multiple users who are running parallel code at the same time. It is also important to note that for a swarm or large user grid platform, provisions of robust grid machine resources is necessary in order to avoid freeze or overload.

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